

Review: What can happen with the Hoare-Calculus

- Hoare Triples can be :
 - not provable (counter-example)
 - provable, but for trivial reasons
 - non termination of the program
 - precondition false (falseE) or equivalent
 - provable for interesting reasons

Exercice 1

- Task 1 : Variante :

$$\vdash \{x \leq 0\} y := y+2 \{y \leq 2\}''$$

- contre-exemple : $y = 5$

- Task 1 as such :

$$\frac{x \leq 0 \longrightarrow y \leq 2 [y \mapsto x+2] \quad \frac{\vdash \{y \leq 2 [y \mapsto x+2]\} y := x+2 \{y \leq 2\} \quad \text{affect} \quad y \leq 2 \longrightarrow y \leq 2}{\text{conseq}}}{\vdash \{x \leq 0\} y := x+2 \{y \leq 2\}}$$

Side calculation :

$$x \leq 0 \longrightarrow y \leq 2 [y \mapsto x+2]$$

$$\equiv x \leq 0 \longrightarrow x+2 \leq 2$$

$$\equiv \text{True}$$

Exercise 1

- Task 2

$$\vdash \{x \leq 0\} y := y+2 \{y \leq 2\}$$

$$\frac{x \leq 0 \longrightarrow x < 0 [x \mapsto x-1] \quad \frac{\vdash \{x < 0 [x \mapsto x-1]\} x := x-1 \{x < 0\} \quad \text{affect} \quad x < 0 \longrightarrow x < 0}{\text{conseq}}}{\vdash \{x \leq 0\} x := x-1 \{x < 0\}}$$

Side Calculations :

$$x \leq 0 \longrightarrow x < 0 [x \mapsto x-1]$$

$$\equiv x \leq 0 \longrightarrow x-1 < 0$$

$$\equiv x \leq 0 \longrightarrow x < 1$$

$$\equiv \text{True}$$

Exercise 1

- Task 3

- Proposition : $I \equiv x \geq -1$

Side Calculations :

$$\begin{aligned} & x \geq -1 [x \mapsto x-1] \\ \equiv & x-1 \geq -1 \equiv x \geq 0 \equiv x \geq -1 \wedge x \geq 0 \end{aligned}$$

$$\frac{x \geq 0 \longrightarrow I \quad \frac{\frac{\frac{}{\vdash \{I \wedge x \geq 0\} x := x-1 \{I\}}{\text{affect}}}{\vdash \{I\} \text{WHILE } x \geq 0 \text{ DO } x := x-1 \{I \wedge x < 0\}} \text{while}}{\vdash \{I \wedge x < 0 \longrightarrow x = -1\}} \text{conseq}}{\vdash \{x \geq 0\} \text{WHILE } x \geq 0 \text{ DO } x := x-1 \{x = -1\}}$$

Exercice 1

Task 4

Prog \equiv $a := a + b; b := a - 2*b; a := a * b$

- Pre $\equiv a = x \wedge b = y$

- Post $\equiv a = x^2 - y^2$

4. On applique deux fois la règle de séquence, et on va appliquer la règle de l'affectation de droite à gauche pour trouver les propriétés intermédiaires R et Q . Puis on devra montrer que $\vdash \{Pre\} a:=a+b \{Q\}$ est valide avec la propriété Q qu'on aura trouvée.

$$\frac{\frac{\frac{?}{\vdash \{Pre\} a:=a+b \{Q\}}}{\vdash \{a = x \wedge b = y\} a:=a+b; b:=a-2*b \{R\}} \text{seq} \quad \frac{\frac{}{\vdash \{Q\} b:=a-2*b \{R\}}}{\vdash \{R\} a:=a*b \{Post\}} \text{aff}}{\vdash \{Pre\} Prog \{Post\}} \text{seq}$$

Exercice 1

- Task 4

Avec $Pre \Leftrightarrow (a = x \wedge b = y)$, $Post \Leftrightarrow (a = x^2 - y^2)$. On a :

$$\begin{aligned} R &\Leftrightarrow Post[a \mapsto a * b] \Leftrightarrow (a * b = x^2 - y^2) \\ Q &\Leftrightarrow R[b \mapsto a - 2 * b] \Leftrightarrow (a^2 - 2 * a * b = x^2 - y^2) \end{aligned}$$

On calcule $Q[a \mapsto a + b]$:

$$Q[a \mapsto a + b] \Leftrightarrow ((a + b)^2 - 2 * (a + b) * b = x^2 - y^2) \Leftrightarrow (a^2 - b^2 = x^2 - y^2)$$

Ce n'est pas directement équivalent à $a = x \wedge b = y$ (si la différence des carrés est égale, on peut aussi avoir $a = -x \wedge b = -y$), mais l'implication $(a = x \wedge b = y) \Rightarrow (a^2 - b^2 = x^2 - y^2)$ est vraie. Donc on pose $P' = a^2 - b^2 = x^2 - y^2$ et on applique la règle de conséquence, pour pouvoir ensuite appliquer la règle de l'affectation.

$$\frac{Pre \Rightarrow P' \quad \frac{}{\vdash \{P'\} a := a + b \{Q\}} \text{aff}}{\vdash \{Pre\} a := a + b \{Q\}} \text{cons}$$

Exercice 1

- Task 4

Observation: it is very difficult to construct R, Q and finally P' from left to right ; however, it is perfectly possible to construct it from right to left and to « bridge » Pre to P' via a consequence rule...

$$\vdash \{a = x \wedge b = y\} a := a + b; b := a - 2*b; a := a * b \{a = x^2 - y^2\}$$

Exercise 1

- Task 5

- Proposition Invariant : $I \equiv i = 8$!!!
- Proposition Invariant : $I \equiv \text{True}$

$$\frac{\begin{array}{c} \frac{}{\vdash \{I \wedge i < 5\} \dots \{I\}} \text{falseE} \\ \vdash \{I\} \text{WHILE } i < 5 \text{ DO } \dots \{I \wedge i \geq 5\} \end{array}}{i=8 \rightarrow I \quad \vdash \{I\} \text{WHILE } i < 5 \text{ DO } i := 2*i \{i \geq 5\}} \text{conseq}$$

Exercice 2

- $A \equiv (\max = x \vee \max = y) \wedge \max \geq x \wedge \max \geq y$

- Justification :

$$\begin{aligned}
 & x > y \longrightarrow A[\max \mapsto x] \\
 \equiv & x > y \longrightarrow (\max = x \vee \max = y) \wedge \max \geq x \wedge \max \geq y [\max \mapsto x] \\
 \equiv & x > y \longrightarrow (x = x \vee x = y) \wedge x \geq x \wedge x \geq y \\
 \equiv & \text{true}
 \end{aligned}$$

$$\begin{array}{c}
 \frac{x > y \longrightarrow A[\max \mapsto x] \quad \frac{}{\vdash \{A[\max \mapsto x]\} \max := x \{A\}} \text{affect}}{\vdash \{A[\max \mapsto x]\} \max := x \{A\}} \text{conseq} \quad A \longrightarrow A \\
 \frac{\vdash \{\text{true}\} \text{IF } x > y \text{ THEN } \max := x \text{ ELSE } \max := y \{A\} \quad \dots}{\vdash \{\text{true} \wedge x > y\} \max := x \{(\max = x \vee \max = y) \wedge \max \geq x \wedge \max \geq y\}} \text{if}
 \end{array}$$

Rappel : La Logique Hoare

Calcul de Hoare

$$\frac{}{\vdash \{P\} \text{ SKIP } \{P\}} \text{skip}$$

$$\frac{}{\vdash \{P[x \mapsto \text{exp}]\} x := \text{exp} \{P\}} \text{aff}$$

$$\frac{\vdash \{P \wedge \text{cond}\} \text{ins}_1 \{Q\} \quad \vdash \{P \wedge \neg \text{cond}\} \text{ins}_2 \{Q\}}{\vdash \{P\} \text{ IF } \text{cond} \text{ THEN } \text{ins}_1 \text{ ELSE } \text{ins}_2 \{Q\}} \text{if}$$

$$\frac{\vdash \{P \wedge \text{cond}\} \text{ins} \{P\}}{\vdash \{P\} \text{ WHILE } \text{cond} \text{ DO } \text{ins} \{P \wedge \neg \text{cond}\}} \text{while}$$

$$\frac{P \Rightarrow P' \quad \vdash \{P'\} \text{ins} \{Q'\} \quad Q' \Rightarrow Q}{\vdash \{P\} \text{ins} \{Q\}} \text{cons}$$

$$\frac{}{\vdash \{false\} \text{ins} \{P\}} \text{falseE}$$

$$\frac{\vdash \{P\} \text{ins}_1 \{Q\} \quad \vdash \{Q\} \text{ins}_2 \{R\}}{\vdash \{P\} \text{ins}_1 ; \text{ins}_2 \{R\}} \text{seq}$$