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Département Informatique

Verification and Validation

Part IV : An Introduction to Testing

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Recall: Validation and Verification

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 - Does the system meet the clients requirements ?
 - Will the performance be sufficient ?
 - Will the usability be sufficient ?

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Do we build the system right ? Is it « correct » ?

How to do Validation ?

- Tests and Experiments over Systems
(Integrated artefacts consisting of software and hardware ...)

How to do Verification ?

- **Test and Proof** on the basis of formal specifications (e.g., à la OCL, MOAL, ACSL, ... !) against programs or systems ...

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 - *on the other side – you can't test infinitely, and verification is again 10 times more costly than thoroughly testing !*

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 - ❑ a lot of knowledge over methods, tools, and tool chains ...

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 - **ideally**: testing is complete if a certain criteria, the adequacy criteria is reached.

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 - non-deterministic programs with no control over the non-determinism.

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 - experiment with the real behaviour
 - essentially used for post-hoc analysis and debugging

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- ❑ **reactive sequence**: testing components by sequences of steps, but these sequences represent communication where later parts in the sequence depend on what has been earlier communicated

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- ❑ **both:** (also: grey-box testing).

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we focus on what the program *should* do !!!

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*Read a "Triangle Object" (with three sides of integral type),
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Each length should be strictly positive.

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Give a specification, and develop a test set ...

Functional Unit Test : An Example

The specification in UML/MOAL:

Triangles

a, b, c: Integer

```
- mk(Integer,Integer,Integer):Triangle
- is_Triangle(): {equ (*equilateral*),
                  iso (*isosceles*),
                  arb (*arbitrary*)}
```

Functional Unit Test : An Example

We add the constraints of
the analysis:

```
inv  0 < a  $\wedge$  0 < b  $\wedge$  0 < c
inv  c  $\leq$  a + b  $\wedge$  a  $\leq$  b + c  $\wedge$  b  $\leq$  c + a
```

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operation t.is_Triangle():

```
post t.a=t.b  $\wedge$  t.b=t.c  $\longrightarrow$  result=equ
```

```
post (t.a $\neq$ t.b  $\vee$  t.b $\neq$ t.c)  $\wedge$ 
```

```
(t.a=t.b  $\vee$  t.b=t.c  $\vee$  t.a=t.c)  $\longrightarrow$  result=iso
```

```
post (t.a $\neq$ t.b  $\vee$  t.b $\neq$ t.c  $\vee$  t.a $\neq$ t.c)  $\longrightarrow$  result=arb
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context C::m( $a_1:C_1, \dots, a_n:C_n$ )  
pre : P(self,  $a_1, \dots, a_n$ )  
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```
check_C(); check_C1(); ... ; check_Cn();  
assert(P(self, a1, ..., an));  
result=run_m(self, a1, ..., an);  
assert(Q(self, a1, ..., an, result));
```

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Dynamic (Unit/Sequence/...) Runtime-Tests are:

- ... easy to implement and enforce
- ... work on real data and are extremely helpful for post-hoc crash-analysis, debugging, and forensics.
- Runtime-tests conflict with efficiency
- But: they are NOT particularly useful during development, where we need systematic test-data EARLY.

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 - ... gives clear criteria on the question:
„did we test enough“ ?

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Note that we define $\text{mk}(0,0,0)$ to be invalid, as well as all other invalid triangles ...

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(1, 2, 4), (4, 1, 2), (2, 4, 1) -- $x + y > z$
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- ❑ Would we have to consider negative values?

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- ❑ Ouf, is there a systematic and automatic way to compute all these tests ?
- ❑ Can we avoid hand-written test-scripts ?
Avoid the task to maintain them ?
- ❑ And the question remains:

When did we test „enough“ ?

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(* see semantics of MOAL in Part III *)

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Some Facts:

- From `modifiesOnly({})` follows $\sigma = \sigma'$ hence

$$\text{inv}_{\text{Triangle}}(\sigma) = \text{inv}_{\text{Triangle}}(\sigma')$$

- From $\text{mk}(x,y,z) \neq \text{null}$ (see $\text{pre}_{\text{isTriangle}}$) and from $\text{inv}_{\text{Triangle}}(\sigma)$ and $\text{mk}(x,y,z) \in \text{Triangle}(\sigma)$ follows that:

$$0 < x \wedge 0 < y \wedge 0 < z \wedge x \leq y + z \wedge y \leq x + z \wedge z \leq x + y \quad (\equiv \text{inv})$$

Revision: Boolean Logic + Some Basic Rules

Revision: Boolean Logic + Some Basic Rules

- ❑ $\neg(a \wedge b) = \neg a \vee \neg b$ (* deMorgan1 *)
- ❑ $\neg(a \vee b) = \neg a \wedge \neg b$ (* deMorgan2 *)
- ❑ $a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$
- ❑ $\neg(\neg a) = a$, $a \vee \neg a = T$, $a \wedge \neg a = F$,
- ❑ $a \wedge b = b \wedge a$; $a \vee b = b \vee a$
- ❑ $a \wedge (b \wedge c) = (a \wedge b) \wedge c$
- ❑ $a \vee (b \vee c) = (a \vee b) \vee c$
- ❑ $a \longrightarrow b = (\neg a) \vee b$
- ❑ $(a=b \wedge P(a)) = P(b)$ (* one point rule *)
- ❑ $\text{let } x = E \text{ in } C(x) = C(E)$ (* let elimination *)
- ❑ $\text{if } c \text{ then } C \text{ else } D = (c \wedge C) \vee (\neg c \wedge D) = (c \longrightarrow C) \wedge (\neg c \longrightarrow D)$

Test-Data Generation

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- Recall the test specification:
`mk(x,y,z).isTriangle() = r`

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(* see semantics d'un appel de methopde, in MOAL II, page 22. *)

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Some Facts:

> $arb \neq equ \neq iso$

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> $\text{post}_{\text{isTriangle}}(\text{mk}(x,y,z),r)(\sigma,\sigma)$ can be simplified to:

$$(x=y \wedge y=z \longrightarrow r=\text{equ}) \wedge$$

$$((x \neq y \vee y \neq z) \wedge (x=y \vee y=z \vee x=z) \longrightarrow r=\text{iso}) \wedge$$

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$$((x \neq y \wedge y \neq z \wedge x \neq z) \longrightarrow r=\text{arb})$$

Test-Data Generation

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\Rightarrow (* the discussed facts *)

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\Rightarrow (* the discussed facts *)

$$\text{inv} \wedge \\ (x=y \wedge y=z \longrightarrow r=\text{equ}) \wedge \\ ((x \neq y \vee y \neq z) \wedge (x=y \vee y=z \vee x=z) \longrightarrow r=\text{iso}) \wedge \\ (x \neq y \wedge y \neq z \wedge x \neq z \longrightarrow r=\text{arb})$$

Test-Data Generation

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≡ (* elimination \longrightarrow , deMorgan*)

$$\begin{aligned} & \text{inv} \wedge \\ & (x \neq y \vee y \neq z \vee r=\text{equ}) \wedge \\ & ((x=y \wedge y=z) \vee (x \neq y \wedge y \neq z \wedge x \neq z) \vee r=\text{iso}) \wedge \\ & (x=y \vee y=z \vee x=z \vee r=\text{arb}) \end{aligned}$$

Test-Data Generation

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- This first part of the calculation could be called

PURIFICATION

We eliminate UML, object-orientation, MOAL etcpp and reduce it to the pure logical core ...

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We eliminate UML, object-orientation, MOAL etcpp and reduce it to the pure logical core ...

Now, under which precise conditions do we have

- > $r = \text{iso}$
- > $r = \text{arb}$
- > $r = \text{equ} \text{ ???}$

Test-Data Generation

- This first part of the calculation could be called

PURIFICATION

We eliminate UML, object-orientation, MOAL etcpp and reduce it to the pure logical core ...

Can we transform the spec into the form

➤ $A_1 \wedge \dots \wedge A_i \wedge r = \text{iso}$

➤ $B_1 \wedge \dots \wedge B_k \wedge r = \text{arb}$

➤ $C_1 \wedge \dots \wedge C_l \wedge r = \text{equ}$???

Test-Data Generation

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PURIFICATION

We eliminate UML, object-orientation, MOAL etcpp and reduce it to the pure logical core ...

Can we transform the spec into a

Disjunctive Normal Form (DNF) ?

Excursion

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- Generalized Distribution Laws:

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$$(A_1 \vee A_2) \wedge (B_1 \vee B_2) = (A_1 \wedge (B_1 \vee B_2)) \vee (A_2 \wedge (B_1 \vee B_2))$$

Excursion

- Generalized Distribution Laws:

$$\begin{aligned}(A_1 \vee A_2) \wedge (B_1 \vee B_2) &= (A_1 \wedge (B_1 \vee B_2)) \vee (A_2 \wedge (B_1 \vee B_2)) \\ &= (A_1 \wedge B_1) \vee (A_2 \wedge B_1) \vee (A_1 \wedge B_2) \vee (A_2 \wedge B_2)\end{aligned}$$

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$$(A_1 \vee A_2 \vee A_3) \wedge (B_1 \vee B_2 \vee B_3) \wedge (C_1 \vee C_2 \vee C_3)$$

Excursion

□ Generalized Distribution Laws:

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$$\begin{aligned}(A_1 \vee A_2 \vee A_3) \wedge (B_1 \vee B_2 \vee B_3) \wedge (C_1 \vee C_2 \vee C_3) \\ = \dots\end{aligned}$$

Excursion

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$$\begin{aligned}(A_1 \vee A_2 \vee A_3) \wedge (B_1 \vee B_2 \vee B_3) \wedge (C_1 \vee C_2 \vee C_3) \\ &= \dots \\ &= (A_1 \wedge B_1 \wedge C_1) \vee (A_1 \wedge B_1 \wedge C_2) \vee (A_1 \wedge B_1 \wedge C_3) \vee \\ &\quad (A_2 \wedge B_1 \wedge C_1) \vee (A_2 \wedge B_1 \wedge C_2) \vee (A_2 \wedge B_1 \wedge C_3) \vee \\ &\quad \dots \\ &\quad (A_1 \wedge B_3 \wedge C_3) \vee (A_2 \wedge B_3 \wedge C_3) \vee (A_3 \wedge B_3 \wedge C_3)\end{aligned}$$

Test-Data Generation

Test-Data Generation

- Recall the test specification:

...

≡ $inv \wedge$

$(x \neq y \vee y \neq z \vee r = equ) \wedge$

$(x = y \vee y = z \vee x = z \vee r = arb) \wedge$

$((x = y \wedge y = z) \vee (x \neq y \wedge y \neq z \wedge x \neq z) \vee r = iso)$

≡

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≡

distrib



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$((x = y \wedge y = z) \vee (x \neq y \wedge y \neq z \wedge x \neq z) \vee r = \text{iso})$

\equiv

$\text{inv} \wedge$

$((x \neq y \wedge x = y) \vee (x \neq y \wedge y = z) \vee (x \neq y \wedge x = z) \vee (x \neq y \wedge r = \text{arb})) \vee$

$((y \neq z \wedge x = y) \vee (y \neq z \wedge y = z) \vee (y \neq z \wedge x = z) \vee (y \neq z \wedge r = \text{arb})) \vee$

$((r = \text{equ} \wedge x = y) \vee (r = \text{equ} \wedge y = z) \vee (r = \text{equ} \wedge x = z) \vee (r = \text{equ} \wedge r = \text{arb})) \vee$

$((x = y \wedge y = z) \vee (x \neq y \wedge y \neq z \wedge x \neq z) \vee r = \text{iso})$

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≡ (* elimination contradictions *)

Test-Data Generation

- Recall the test specification:

...

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$(x \neq y \vee y \neq z \vee r = equ) \wedge$

$(x = y \vee y = z \vee x = z \vee r = arb) \wedge$

$((x = y \wedge y = z) \vee (x \neq y \wedge y \neq z \wedge x \neq z) \vee r = iso)$

≡ **(* elimination contradictions *)**

$inv \wedge$

$((x \neq y \wedge x = y) \vee (x \neq y \wedge y = z) \vee (x \neq y \wedge x = z) \vee (x \neq y \wedge r = arb) \vee$

$(y \neq z \wedge x = y) \vee (y \neq z \wedge y = z) \vee (y \neq z \wedge x = z) \vee (y \neq z \wedge r = arb) \vee$

$(r = equ \wedge x = y) \vee (r = equ \wedge y = z) \vee (r = equ \wedge x = z) \vee (r = equ \wedge r = arb)) \vee$

$((x = y \wedge y = z) \vee (x \neq y \wedge y \neq z \wedge x \neq z) \vee r = iso)$

Test-Data Generation

- Recall the test specification:

...

≡ (* elimination contradictions *)

inv \wedge

$$\begin{aligned} & \left((x \neq y \wedge y = z) \vee (x \neq y \wedge x = z) \vee (x \neq y \wedge r = \text{arb}) \vee \right. \\ & \quad (y \neq z \wedge x = y) \vee (y \neq z \wedge x = z) \vee (y \neq z \wedge r = \text{arb}) \vee \\ & \quad \left. (r = \text{equ} \wedge x = y) \vee (r = \text{equ} \wedge y = z) \vee (r = \text{equ} \wedge x = z) \right) \wedge \\ & \left((x = y \wedge y = z) \vee (x \neq y \wedge y \neq z \wedge x \neq z) \vee r = \text{iso} \right) \end{aligned}$$

Test-Data Generation

- ≡ (* generalized distribution 2nd/3rd ((9 * 3 = 27 cases !)*
- inv \wedge
- $$\begin{aligned} & \left((x \neq y \wedge y = z \wedge x = y \wedge y = z) \vee (x \neq y \wedge x = z \wedge \right. \\ & \qquad \qquad \qquad \left. x = y \wedge y = z) \vee (x \neq y \wedge r = \text{arb} \wedge x = y \wedge y = z) \vee \right. \\ & (y \neq z \wedge x = y \wedge x = y \wedge y = z) \vee (y \neq z \wedge x = z \wedge \\ & \qquad \qquad \qquad \left. x = y \wedge y = z) \vee (y \neq z \wedge r = \text{arb} \wedge x = y \wedge y = z) \vee \right. \\ & (r = \text{equ} \wedge x = y \wedge x = y \wedge y = z) \vee (r = \text{equ} \wedge \\ & \qquad \qquad \qquad \left. y = z \wedge x = y \wedge y = z) \vee (r = \text{equ} \wedge x = z \wedge x = y \wedge y = z) \right) \vee \\ & \left((x \neq y \wedge y = z \wedge x \neq y \wedge y \neq z \wedge x \neq z) \vee (x \neq y \wedge x = z \wedge x \neq y \wedge y \neq z \wedge x \neq z) \vee (x \neq y \wedge r = \text{arb} \wedge \right. \\ & \left. x \neq y \wedge y \neq z \wedge x \neq z) \vee (y \neq z \wedge x = y \wedge x \neq y \wedge y \neq z \wedge x \neq z) \vee (y \neq z \wedge x = z \wedge x \neq y \wedge y \neq z \wedge \right. \\ & \left. x \neq z) \vee (y \neq z \wedge r = \text{arb} \wedge x \neq y \wedge y \neq z \wedge x \neq z) \vee (r = \text{equ} \wedge x = y \wedge x \neq y \wedge y \neq z \wedge x \neq z) \vee \right. \\ & \left. r = \text{equ} \wedge y = z \wedge x \neq y \wedge y \neq z \wedge x \neq z) \vee (r = \text{equ} \wedge x = z \wedge x \neq y \wedge y \neq z \wedge x \neq z) \right) \vee \\ & \left((x \neq y \wedge y = z \wedge r = \text{iso}) \vee (x \neq y \wedge x = z \wedge r = \text{iso}) \vee (x \neq y \wedge r = \text{arb} \wedge r = \text{iso}) \vee \right. \\ & \left. (y \neq z \wedge x = y \wedge r = \text{iso}) \vee (y \neq z \wedge x = z \wedge r = \text{iso}) \vee (y \neq z \wedge r = \text{arb} \wedge r = \text{iso}) \vee \right. \\ & \left. (r = \text{equ} \wedge x = y \wedge r = \text{iso}) \vee (r = \text{equ} \wedge y = z \wedge r = \text{iso}) \vee (r = \text{equ} \wedge x = z \wedge r = \text{iso}) \right) \end{aligned}$$

Test-Data Generation

- ≡ (* elimination of the contradictions and redundancies *)

inv \wedge

$$\begin{aligned}
 & \left((x \neq y \wedge y = z \wedge x = y \wedge y = z) \vee (x \neq y \wedge x = z \wedge \right. \\
 & \qquad \qquad \qquad \left. x = y \wedge y = z) \vee (x \neq y \wedge r = \text{arb} \wedge x = y \wedge y = z) \vee \right. \\
 & \left. (y \neq z \wedge x = y \wedge x = y \wedge y = z) \vee (y \neq z \wedge x = z \wedge \right. \\
 & \qquad \qquad \qquad \left. x = y \wedge y = z) \vee (y \neq z \wedge r = \text{arb} \wedge x = y \wedge y = z) \vee \right. \\
 & \left. (r = \text{equ} \wedge x = y \wedge x = y \wedge y = z) \vee \underline{(r = \text{equ} \wedge \right.} \\
 & \qquad \qquad \qquad \left. \underline{y = z \wedge x = y \wedge y = z)} \vee \underline{(r = \text{equ} \wedge x = z \wedge x = y \wedge y = z)} \right) \vee \\
 & \left((x \neq y \wedge y = z \wedge x \neq y \wedge y \neq z \wedge x \neq z) \vee (x \neq y \wedge x = z \wedge x \neq y \wedge y \neq z \wedge x \neq z) \vee (x \neq y \wedge r = \text{arb} \wedge \right. \\
 & \left. \underline{x \neq y \wedge y \neq z \wedge x \neq z}) \vee (y \neq z \wedge x = y \wedge x \neq y \wedge y \neq z \wedge x \neq z) \vee (y \neq z \wedge x = z \wedge x \neq y \wedge y \neq z \wedge \right. \\
 & \left. \underline{x \neq z}) \vee \underline{(y \neq z \wedge r = \text{arb} \wedge x \neq y \wedge y \neq z \wedge x \neq z)} \vee (r = \text{equ} \wedge x = y \wedge x \neq y \wedge y \neq z \wedge x \neq z) \vee \right. \\
 & \left. r = \text{equ} \wedge y = z \wedge x \neq y \wedge y \neq z \wedge x \neq z) \vee (r = \text{equ} \wedge x = z \wedge x \neq y \wedge y \neq z \wedge \underline{x \neq z}) \right) \vee \\
 & \left((x \neq y \wedge y = z \wedge r = \text{iso}) \vee (x \neq y \wedge x = z \wedge r = \text{iso}) \vee (x \neq y \wedge r = \text{arb} \wedge r = \text{iso}) \vee \right. \\
 & \left. (y \neq z \wedge x = y \wedge r = \text{iso}) \vee (y \neq z \wedge x = z \wedge r = \text{iso}) \vee (y \neq z \wedge r = \text{arb} \wedge r = \text{iso}) \vee \right. \\
 & \left. (r = \text{equ} \wedge x = y \wedge r = \text{iso}) \vee (r = \text{equ} \wedge y = z \wedge r = \text{iso}) \vee (r = \text{equ} \wedge x = z \wedge r = \text{iso}) \right)
 \end{aligned}$$

Test-Data Generation

Test-Data Generation

□ ≡ (* cleanup, distribution *)

(inv \wedge x=y \wedge x=y \wedge y=z \wedge r=equ) V (1)

(inv \wedge x \neq y \wedge y \neq z \wedge x \neq z \wedge r=arb) V (2)

(inv \wedge x \neq y \wedge y=z \wedge r=iso) V (3)

(inv \wedge x \neq y \wedge x=z \wedge r=iso) V (4)

(inv \wedge y \neq z \wedge x=y \wedge r=iso) V (5)

(inv \wedge y \neq z \wedge x=z \wedge r=iso) (6)

Test-Data Generation

□ ≡ (* cleanup, distribution *)

$$(\text{inv} \wedge x=y \wedge x=y \wedge y=z \wedge r=\text{equ}) \quad \mathbf{V} \quad (1)$$

$$(\text{inv} \wedge x \neq y \wedge y \neq z \wedge x \neq z \wedge r=\text{arb}) \quad \mathbf{V} \quad (2)$$

$$(\text{inv} \wedge x \neq y \wedge y=z \wedge r=\text{iso}) \quad \mathbf{V} \quad (3)$$

$$(\text{inv} \wedge x \neq y \wedge x=z \wedge r=\text{iso}) \quad \mathbf{V} \quad (4)$$

$$(\text{inv} \wedge y \neq z \wedge x=y \wedge r=\text{iso}) \quad \mathbf{V} \quad (5)$$

$$(\text{inv} \wedge y \neq z \wedge x=z \wedge r=\text{iso}) \quad \mathbf{V} \quad (6)$$

□ Test-Case-Construction by DNF Method

Test-Data Generation

□ ≡ (* cleanup, distribution *)

$$(inv \wedge x=y \wedge x=y \wedge y=z \wedge r=equ) \vee \quad (1)$$

$$(inv \wedge x \neq y \wedge y \neq z \wedge x \neq z \wedge r=arb) \vee \quad (2)$$

$$(inv \wedge x \neq y \wedge y=z \wedge r=iso) \vee \quad (3)$$

$$(inv \wedge x \neq y \wedge x=z \wedge r=iso) \vee \quad (4)$$

$$(inv \wedge y \neq z \wedge x=y \wedge r=iso) \vee \quad (5)$$

$$(inv \wedge y \neq z \wedge x=z \wedge r=iso) \quad (6)$$

□ Test-Case-Construction by DNF Method

yields six abstract test cases

relating input x y z to output r

Test-Data Generation

□ ≡ (* cleanup, distribution *)

(inv \wedge x=y \wedge x=y \wedge y=z \wedge r=equ) \vee (1)

(inv \wedge x \neq y \wedge y \neq z \wedge x \neq z \wedge r=arb) \vee (2)

(inv \wedge x \neq y \wedge y=z \wedge r=iso) \vee (3)

(inv \wedge x \neq y \wedge x=z \wedge r=iso) \vee (4)

(inv \wedge y \neq z \wedge x=y \wedge r=iso) \vee (5)

(inv \wedge y \neq z \wedge x=z \wedge r=iso) (6)

□ Test-Case-Construction by DNF Method

yields six abstract test cases

relating input x y z to output r

□ Note: In general, output r is not necessarily **uniquely defined** as in our example ...

The spec can be non-deterministic admitting several results.

Test-Data Generation

Test-Data Generation

- Test-Data-Selection:

For each abstract test-case, we construct one concrete test, by choosing values that make the abstract test case true (« that satisfies the abstract test case »)

Test-Data Generation

□ Test-Data-Selection:

For each abstract test-case, we construct one concrete test, by choosing values that make the abstract test case true (« that satisfies the abstract test case »)

case	x	y	z	result
(1)	3	3	3	equ
(2)	3	4	6	arb
(3)	4	5	5	iso
(4)	5	4	5	iso
(5)	5	5	4	iso
(6)	4	3	4	iso

Test-Data Generation

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 - ❑ ... provided that it is not redundant (“A=True”)

Test-Data Generation

- ❑ Intuitively, what does it mean that we “covered” the DNF by tests
 - ❑ Any basic predicate (“literal”) has been used at least one time
 - ❑ ... provided it is not contradictory (“A=False”)
 - ❑ ... provided that it is not redundant (“A=True”)
 - ❑ ... provided it is not implied by another literal, i.e. it is subsumed (“B \longrightarrow A”)

Test-Data Generation

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- A First Summary on the Test-Generation Method:

Test-Data Generation

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 - PHASE I: Stripping the Domain-Language (UML-MOAL) away, "purification"

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DNF - coverage of the Spec; for each abstract test-case one concrete test-input is constructed.
(ISO/IEC/IEEE 29119 calls this: Equivalence class testing)

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 - COVERAGE CRITERION:
DNF - coverage of the Spec; for each abstract test-case one concrete test-input is constructed.
(ISO/IEC/IEEE 29119 calls this: Equivalence class testing)
- ❑ Remark: During Coding phase, when the Spec does not change, the test-data-selection can be repeated easily creating always different test sets ...

Test-Data Generation

□ Variants:

- Alternative to PHASE II (DNF construction):
Predicate Abstraction and Tableaux-Exploration.

Reconsider the (purified) specification:

$$\begin{aligned} & \text{inv} \wedge \\ & (x=y \wedge y=z \longrightarrow r=\text{equ}) \wedge \\ & ((x \neq y \vee y \neq z) \wedge (x=y \vee y=z \vee x=z) \longrightarrow r=\text{iso}) \wedge \\ & (x \neq y \wedge y \neq z \wedge x \neq z \longrightarrow r=\text{arb}) \end{aligned}$$

It is possible to abstract this spec to a fairly small number of „base predicates“ ... They should be logically independent and not contain the output variable...

Test-Data Generation

□ Variants:

- Alternative to PHASE II (DNF construction):
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Reconsider the (purified) specification:

$$\begin{aligned} & \text{inv} \wedge \\ & (A \wedge B \longrightarrow r=\text{equ}) \wedge \\ & ((\neg A \vee \neg B) \wedge (A \vee B \vee C) \longrightarrow r=\text{iso}) \wedge \\ & (\neg A \wedge \neg B \wedge \neg C \longrightarrow r=\text{arb}) \end{aligned}$$

where $A \mapsto x=y$, $B \mapsto y=z$, $C \mapsto x=z$

(actually: A and B imply C)

Test-Data Generation

□ Variants:

➤ ... Now we can construct a tableau and get by simplification:

case	A	B	C	spec reduces to
(1)	T	T	T	• r=equ
(2)	T	T	F	• r=equ (!!!)
(3)	T	F	T	• r=iso
(4)	T	F	F	• r=iso
(5)	F	T	T	• r=iso
(6)	F	T	F	• r=iso
(7)	F	F	T	• r=iso
(8)	F	F	F	• r=arb

Test-Data Generation

▣ Variants:

- PHASE III: Borderline analysis.

Principle: we replace in our DNF inequalities by „the closest values that make the spec true“

$$x \neq y \quad \mapsto \quad x = y + 1 \quad \vee \quad x = y - 1$$

$$x \leq y \quad \mapsto \quad x = y \quad \vee \quad x < y$$

$$x < y \quad \mapsto \quad x = y - 1 \quad \text{etc.}$$

- ... and recompute the DNF. In general, this gives a much finer mesh ...

Test-Data Generation

□ Variants:

- PHASE I: Test for exceptional behaviour.

We negate the precondition and to DNF generation on the precondition only.

Test objectives could be:

- should raise an exception if public
- should not diverge

Test-Data Generation

- How to handle Recursion ?

Test-Data Generation

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Test-Data Generation

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- at the level
of data

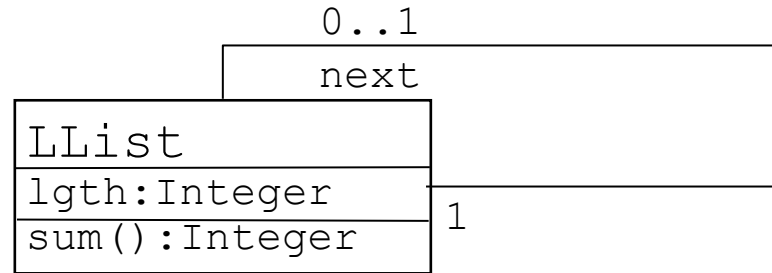
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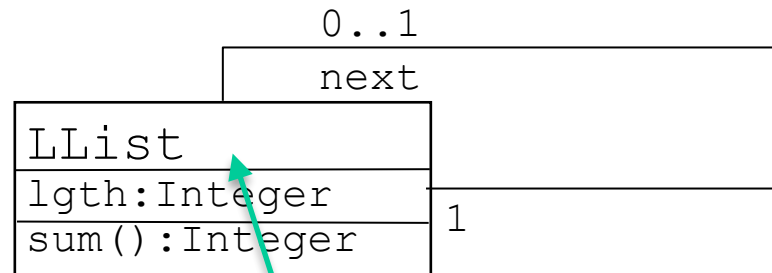
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invariant:
 $inv_{LList} \equiv \forall node \in LList. \quad \begin{array}{l} node.lgth = \text{if } node.next = \text{null} \\ \text{then } 1 \\ \text{else } next.lgth + 1 \end{array}$

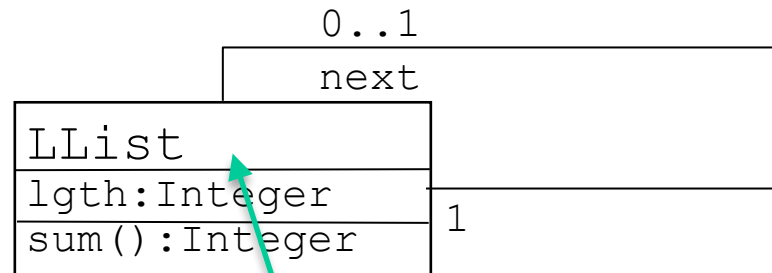
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 $inv_{LList} \equiv \forall node \in LList. \quad node.lgth = \begin{cases} 1 & \text{if } node.next = null \\ next.lgth + 1 & \text{else} \end{cases}$

Note that this excludes cyclic lists !!!

Test-Data Generation

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In UML/MOAL, recursion occurs (at least) at two points:

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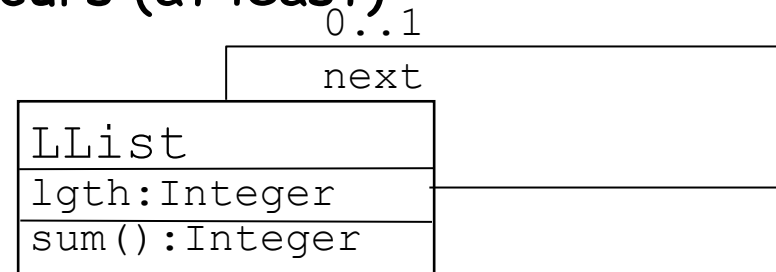
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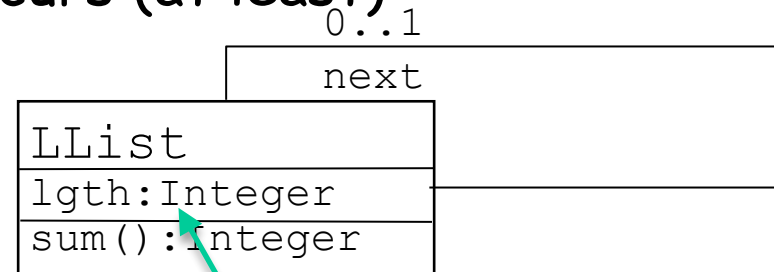
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```
query contract (modifiesOnly({})):
definition  $pre_{sum}(l) \equiv True$ 
definition  $post_{sum}(l, res) \equiv res = if\ l.next = null\ then\ l.lgth$ 
                                      $else\ l.lgth + l.next.sum()$ 
definition  $sum(l) \equiv arb\{r | pre_{sum}(l) \wedge post_{sum}(l, r)\}$ 
```

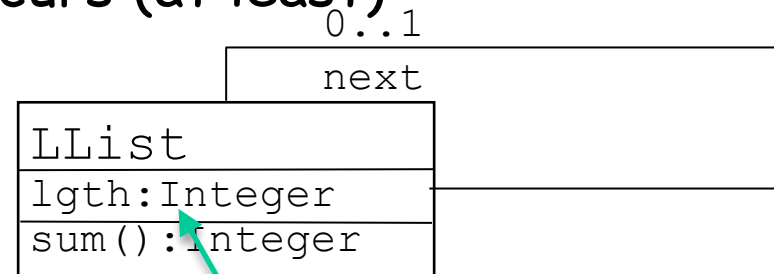
Test-Data Generation

How to handle Recursion ?

In UML/MOAL, recursion occurs (at least)

at two points:

- at the level of operations (post-conds may contain calls ...)



```
query contract (modifiesOnly({})):
definition presum(l) ≡ True
definition postsum(l, res) ≡ res = if l.next = null then l.lgth
                                else l.lgth + l.next.sum()
definition sum(l) ≡ arb{r | presum(l) ∧ postsum(l, r)}
```

Note that $\text{arb}(S)$ gives an arbitrary member of S : $\text{arb}(S) \in S$. Since from $x = \text{arb}(\{y\})$ follows $x = y$; thus $\text{sum}(l)$ is (uniquely) defined.

Test-Data Generation

- Prerequisite: We present the invariant as recursive predicate.

definition $\text{inv}_{\text{LList_Core}}^n \sigma \equiv (\text{n.lgth}(\sigma) = \text{if } \text{n.next}(\sigma) = \text{null} \text{ then } 1$
 $\text{else } \text{n.next.lgth}(\sigma) + 1)$

we have:

$$\text{inv}_{\text{LList}}(\sigma) = \forall n \in \text{LList}(\sigma). \text{inv}_{\text{LList_Core}}^n \sigma$$

and

$$\text{inv}_{\text{LList_Core}}(n)(\sigma) = (\text{if } \text{n.next}(\sigma) = \text{null} \text{ then } \text{n.lgth}(\sigma) = 1$$
$$\text{else } \text{n.lgth}(\sigma) = \text{n.next.lgth}(\sigma) + 1$$
$$\wedge \text{n.next}(\sigma) \in \text{LList}(\sigma)$$
$$\wedge \text{inv}_{\text{LList_Core}}(\text{n.next})(\sigma))$$

Furthermore we have:

$$\text{sum}(l)(\sigma', \sigma) = \text{if } \text{l.next}(\sigma) = \text{null} \text{ then } \text{l.lgth}(\sigma)$$
$$\text{else } \text{l.lgth}(\sigma) + \text{sum}(\text{l.next})(\sigma', \sigma)$$

We have $\sigma' = \sigma$ (why?). We will again apply (σ', σ) - convention.

Test-Data Generation

- Consider the test specification:

$$X.\text{sum}() \equiv Y \quad (\text{for some } X \in \text{LList, i.e. } X \neq \text{null})$$

$$\equiv \text{inv}_{\text{LList}}(X) \wedge \text{pre}_{\text{sum}}(X) \wedge \text{post}_{\text{sum}}(X, Y)$$

where:

$$\text{pre}_{\text{sum}}(X) \equiv \text{true}$$

$$\begin{aligned} \text{post}_{\text{sum}}(X, Y) &\equiv (\text{if } X.\text{next} = \text{null} \text{ then } Y = X.\text{lgth} \\ &\quad \text{else } Y = X.\text{lgth} + \text{sum}(X.\text{next})) \\ &\equiv (X.\text{next} = \text{null} \wedge Y = X.\text{lgth}) \\ &\quad \vee (X.\text{next} \neq \text{null} \wedge Y = X.\text{lgth} + \text{sum}(X.\text{next})) \end{aligned}$$

Test-Data Generation

Test-Data Generation

- DNF computation yields already the test cases:

$$X.\text{sum}() \equiv Y \quad (\text{for some } X \in \text{LList, i.e. } X \neq \text{null})$$

$$\Rightarrow \text{inv}_{\text{LList_Core}}(X) \wedge \text{post}_{\text{sum}}(X, Y)$$

$$\equiv (\text{if } X.\text{next}=\text{null} \text{ then } X.\text{lgth} = 1 \\ \text{else } X.\text{lgth} = X.\text{next}.\text{lgth}+1 \wedge X.\text{next} \in \text{LList} \wedge \text{inv}_{\text{LList_Core}}(X.\text{next})) \wedge \\ (\text{if } X.\text{next} = \text{null} \text{ then } Y = X.\text{lgth} \\ \text{else } Y = X.\text{lgth} + \text{sum}(X.\text{next}))$$

$$\equiv (\text{if } c \text{ then } C \text{ else } D \text{ elim, DNF})$$

$$(X.\text{next}=\text{null} \wedge X.\text{lgth}=1 \wedge Y = X.\text{lgth}) \\ \vee (X.\text{next} \neq \text{null} \wedge X.\text{lgth} = X.\text{next}.\text{lgth}+1 \\ \wedge X.\text{next} \in \text{LList} \wedge \text{inv}_{\text{LList_Core}}(X.\text{next}) \\ \wedge Y = X.\text{lgth} + \text{sum}(X.\text{next}))$$

Test-Data Generation

- DNF computation yields already the test cases:

$X.sum() \equiv Y$ (for some $X \in LList$, i.e. $X \neq null$)

$\Rightarrow inv_{LList_Core}(X) \wedge post_{sum}(X, Y)$

\equiv (if $X.next = null$ then $X.lgth = 1$
else $X.lgth = X.next.lgth + 1 \wedge X.next \in LList \wedge inv_{LList_Core}(X.next)$) \wedge
(if $X.next = null$ then $Y = X.lgth$
else $Y = X.lgth + sum(X.next)$)

\equiv (if c then C else D elim, DNF)

$(X.next = null \wedge X.lgth = 1 \wedge Y = X.lgth)$

$\vee (X.next \neq null \wedge X.lgth = X.next.lgth + 1$
 $\wedge X.next \in LList \wedge inv_{LList_Core}(X.next)$
 $\wedge Y = X.lgth + sum(X.next))$

New
Test-
Case!!

Test-Data Generation

Test-Data Generation

- ❑ Intermediate Summary: test-cases known so far ?

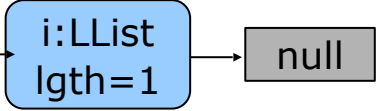
Test-Data Generation

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X	Y
	1
...	...
...	...

Test-Data Generation

- Intermediate Summary: test-cases known so far ?

X	Y
 <p>A diagram showing a blue rounded rectangle containing the text 'i:LList' and 'lgth=1'. An arrow points from the left side of this rectangle to a grey rectangle containing the text 'null'.</p>	1
...	...
...	...

Test-Data Generation

- Prerequisite: We present the invariant as recursive predicate.

$$\text{inv}_{\text{LList_Core}}(n) = (\text{if } n.\text{next}=\text{null} \text{ then } n.\text{lgth} = 1 \\ \text{else } n.\text{lgth} = n.\text{next}.\text{lgth} + 1 \\ \wedge n.\text{next} \in \text{LList} \wedge \text{inv}_{\text{LList_Core}}(n.\text{next}))$$

- $\text{sum}(l) = \text{if } l.\text{next}=\text{null} \text{ then } l.\text{lgth} \\ \text{else } l.\text{lgth} + \text{sum}(l.\text{next})$

$$\text{sum}(l) = \text{if } X.\text{next}.\text{next}=\text{null} \text{ then } X.\text{next}.\text{lgth} \\ \text{else } X.\text{next}.\text{lgth} + \text{sum}(X.\text{next}.\text{next})$$

Test-Data Generation

- DNF computation yields already the test cases:

$$X.\text{sum}() \equiv Y \quad (\text{for some } X \in \text{LList, i.e. } X \neq \text{null})$$

$\implies \dots \equiv \dots$

$\equiv (\text{unfolding sum and } \text{inv}_{\text{LList_Core}})$

$(X.\text{next}=\text{null} \wedge X.\text{lgth}=1 \wedge Y = X.\text{lgth})$

$\vee (X.\text{next} \neq \text{null} \wedge X.\text{lgth}=X.\text{next}.\text{lgth}+1 \wedge X.\text{next} \in \text{LList}$

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Test-Data Generation

- DNF computation yields already the test cases:

$$X.\text{sum}() \equiv Y \quad (\text{for some } X \in \text{LList, i.e. } X \neq \text{null})$$

$\Rightarrow \dots \equiv \dots$

\equiv (DNF partial)

$$\begin{aligned} & (X.\text{next}=\text{null} \wedge X.\text{lgth}=1 \wedge Y = X.\text{lgth}) \\ & \vee (X.\text{next} \neq \text{null} \wedge X.\text{lgth}=X.\text{next}.\text{lgth}+1 \wedge X.\text{next} \in \text{LList} \\ & \wedge ((X.\text{next}.\text{next}=\text{null} \wedge X.\text{next}.\text{lgth} = 1 \wedge Y = X.\text{lgth}+X.\text{next}.\text{lgth}) \\ & \vee (X.\text{next}.\text{next} \neq \text{null} \wedge X.\text{next}.\text{lgth}=X.\text{next}.\text{next}.\text{lgth}+1 \\ & \quad \wedge X.\text{next}.\text{next} \in \text{LList} \wedge \text{inv}_{\text{LList_Core}}(X.\text{next}.\text{next}) \\ & \quad \wedge Y = X.\text{lgth}+ X.\text{next}.\text{lgth} + \text{sum}(X.\text{next}.\text{next})) \\ &) \end{aligned}$$

Test-Data Generation

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- DNF computation yields already the test cases:

$X.sum() \equiv Y$ (for some $X \in LList$, i.e. $X \neq null$)

$\Rightarrow \dots \equiv \dots$

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$(X.next=null \wedge X.lgth=1 \wedge Y = X.lgth)$

$\vee (X.next \neq null \wedge X.lgth=X.next.lgth+1 \wedge X.next \in LList$
 $\wedge X.next.next=null \wedge X.next.lgth=1 \wedge Y = X.lgth+X.next.lgth))$

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Test-Data Generation

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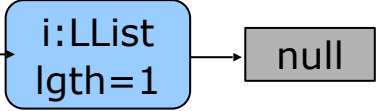
Test-Data Generation

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X	Y
	1
...	2
...	...

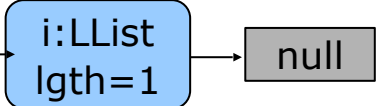
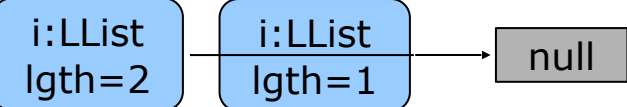
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...	...

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 - simplify DNF
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- until we are satisfied, i.e. have „enough“ test cases ...
- **Select test-data:** constraint resolution of test cases.

Test-Data Generation

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- **Observation:** "all other cases" ...
were represented by the clauses still
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Test-Data Generation

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- **Logically:** we used a **regularity hypothesis**, i.e ...

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Test-Data Generation

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were represented by the clauses still
containing recursive predicates.
- **Logically:** we used a **regularity hypothesis**, i.e ...

$$\begin{aligned}(\forall X. |X| < k \Rightarrow X.\text{sum}() \equiv Y) \\ \Rightarrow (\forall X. X.\text{sum}() \equiv Y)\end{aligned}$$

where we choose as "complexity measure" $|X|$
just $X.\text{lgth}$ and k (the number of unfoldings)
was 2 ...

Test-Data Generation

- Coverage Criterion for recursive specification:

$$\text{DNF}_k$$

For all data up to complexity k , we constructed abstract test-cases and generated a test.

In our example, the “complexity measure” is just the length of the LLists.

Test-Data Generation

- ❑ What are the alternatives to symbolic test-case generation ?

Must this really be so complicated ???

Well, think about the probability to "guess" input with a complex invariant or precondition, if you use "blind" random-generation of input...

Test-Data Generation

Test-Data Generation

- Summary

Test-Data Generation

- Summary

- We have (sketched) a symbolic Test-Case Generation Procedure for UML/MOAL Specifications

Test-Data Generation

□ Summary

- We have (sketched) a symbolic Test-Case Generation Procedure for UML/MOAL Specifications
- It takes into account:

Test-Data Generation

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 - data invariants (recursive predicates)
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- The process can be tool-supported (HOL-TestGen)
- The process is intended for automation.

Test-Data Generation

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- Adequacy:
Pick for each test-case (a conjunct in the DNF_k)
one test, i.e. one substitution for the free
variables satisfying the test-case !