

# TD 9

## Exercice 1

1. Si possible, dérivez les triplets de Hoare suivants en utilisant les règles d'inférence introduites dans le cours (ces règles sont rappelées au dos de cette page). On suppose l'arithmétique des nombres entiers.

- (a)  $\vdash \{x > 0\} x := x - 1 \{0 \leq x\}$
- (b)  $\vdash \{0 \leq x \wedge x \bmod 5 > 5\} x := x * x \{x \bmod 2 = 1\}$
- (c)  $\vdash \{x \leq 0\} \text{WHILE } x < 0 \text{ DO } x := x + 1 \{x = 0\}$
- (d)  $\vdash \{S * Y^P = X^N\} P := P - 1; S := S * Y \{S * Y^P = X^N\}$
- (e)  $\vdash \{x \leq -2\} \text{WHILE } 0 < x * x \text{ DO } x := x + 1 \{x = 0\}$

**Figure 1: Preuves Simples**

(a)

$\frac{x > 0 \rightarrow (0 \leq x [x \mapsto x - 1]) \quad \frac{}{\vdash \{0 \leq x [x \mapsto x - 1]\} x := x - 1 \{0 \leq x\}} \text{(aff)} \quad 0 \leq x \rightarrow 0 \leq x}{\vdash \{x > 0\} x := x - 1 \{0 \leq x\}}$
<pre> ` {x &gt; 0} x := x - 1 {0 ≤ x}  x &gt; 0 -&gt; (0 ≤ x [x -&gt; x - 1]) equiv x &gt; 0 -&gt; x - 1 equiv x &gt; 0 -&gt; x &gt;= 1 equiv True         </pre>

(b)

$$\frac{\text{falseE}}{\vdash \{0 \leq x \wedge x \bmod 5 > 5\} x := x * x \{x \bmod 2 = 1\}}$$


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$$\vdash \{0 \leq x \wedge x \bmod 5 > 5\} x := x * x \{x \bmod 2 = 1\}$$

(c)

$$\frac{\frac{\text{B } \{I \wedge x < 0 \rightarrow I[x \rightarrow x + 1]\} \quad \frac{\{I[x := x + 1]\} \quad x := x + 1 \{I\} \text{ (aff)} \quad I \rightarrow I}{\vdash \{I \wedge x < 0\} \quad x := x + 1 \{I\} \text{ (cons)}}}{\text{A } \{x \leq 0 \rightarrow I\} \quad \frac{\vdash \{I\} \text{ WHILE } \dots \{I \wedge \text{not}(x < 0)\} \text{ (while)} \quad \{I \wedge \text{not}(x < 0) \rightarrow x = 0\}}{\vdash \{x \leq 0\} \text{ WHILE } x < 0 \text{ DO } x := x + 1 \{x = 0\} \text{ (cons)}}$$

On cherche I de sorte que :

- A)  $x \leq 0 \rightarrow I$
- B)  $I \wedge x < 0 \rightarrow I[x \rightarrow x + 1]$
- C)  $I \wedge \text{not}(x < 0) \rightarrow x = 0$

side proof:

$x \leq 0 \rightarrow x \leq 0$  Inv établi

$x \leq 0 \wedge x < 0 \rightarrow (x \leq 0)[x \rightarrow x + 1]$

equiv  $x < 0 \rightarrow x + 1 \leq 0$

equiv  $x < 0 \rightarrow x \leq -1$

Inv establishes post:

$x \leq 0 \wedge \text{not}(x < 0) \rightarrow x = 0$

equiv  $x \leq 0 \wedge x \geq 0 \rightarrow x = 0$

equiv True

(d)

$\frac{}{\vdash \{Q[P \mapsto P-1]\} P:=P-1 \{Q\}} \text{ assign}$	$\frac{}{\vdash \{Q\} s:=S * Y \{S * Y^P = X^P\}} \text{ assign}$
$\vdash \{S * Y^P = X^N\} P:=P-1; S:=S*Y \{S * Y^P = X^N\} \quad \text{(sequence)}$	
<p>Q equiv <math>(S * Y^P = X^N)[S \mapsto S * Y]</math>  equiv <math>S * Y * Y^P = X^N</math>    <math>\Leftarrow</math>  equiv <math>S * Y^{(P+1)} = X^N</math>    <math>\Leftarrow</math></p>	<p><math>Q[P \mapsto P-1]</math>  equiv <math>S * Y^{(P+1)} = X^N [P \mapsto P-1]</math>  equiv <math>S * Y^{((P-1)+1)} = X^N</math>  equiv <math>(S * Y^P = X) = X^N</math></p> <p>ce qui correspond a la precondition globale....</p>

e)

	$\frac{}{\vdash \{I \wedge 0 < x * x \mapsto I[x \mapsto x+1]\} x:=x+1 \{I\}} \text{ (aff)}$	
	$\frac{}{\vdash \{I \wedge 0 < x * x\} x:=x+1 \{I\}} \text{ (cons)}$	
$x \leq -2 \mapsto I$	$\frac{}{\vdash \{I\} \text{ WHILE } 0 < x * x \text{ DO } \dots \{I \wedge x * x \geq 0\}} \text{ (while)}$	$I \wedge x * x \geq 0 \mapsto x=0$
$\vdash \{x \leq -2\} \text{ WHILE } 0 < x * x \text{ DO } x:=x+1 \{x=0\}$		
<p>Hypothesis: <math>I \equiv x \leq 0</math></p>	<p>Check A : <math>x \leq -2 \mapsto I \equiv x \leq -2 \mapsto x \leq 0 \equiv \text{True}</math></p> <p>Check B : <math>I \wedge 0 &lt; x * x \mapsto I[x \mapsto x+1]</math>  <math>\equiv x \leq 0 \wedge 0 &lt; x * x \mapsto x+1 \leq 0</math>  <math>\equiv x \leq 0 \wedge 0 &lt; x * x \wedge 0 \neq x \mapsto x \leq -1</math>  <math>\equiv \text{True}</math></p> <p>Check C : <math>I \wedge x * x \geq 0 \mapsto x=0 \equiv x \leq 0 \wedge x * x \geq 0 \mapsto x=0 \equiv \text{True}</math></p>	